

KECS-301 DISCRETE STRUCTURES & THEORY OF LOGIC (DSTL)

① Literal meaning.

DISCRETE :- \rightarrow Noncontinuous. [^{Ex:} Discrete sets include finite and countable sets]
collective

MATHEMATICAL STRUCTURES/MATHEMATICS :- form, [^{Ex:} structures that are discrete are Combinations, graphs, and logical statements.]

THEORY OF LOGIC :- \rightarrow The sensible/reasonable way of thinking.

② why Study. [an idea or set of ideas that tries to explain something]

\Rightarrow It is a vital prerequisite to learning algorithms, as it covers probabilities, trees, graphs, logic, mathematical thinking and much more. It simply explains them, so once you get these basic topics, it is easier to dig into algorithms.

\Rightarrow The problem-solving techniques honed in a discrete mathematics are necessary for writing complicated software. Students who are successful in discrete mathematics will be able to generalize from a single instance of a problem to an entire class of problems, and to identify and abstract patterns from data.

→

"It's a building block for logical thinking."

③ Real life example:- "Wiring a computer network using the least amount of cable is a minimum-weight spanning tree problem."

④ Discrete vs Continuous

\Rightarrow Discrete data is the type of data that has clear spaces between values. Continuous data is data that falls in a constant sequence.
 \Rightarrow Discrete data is countable, while continuous - measurable.

⑤ Ex:- Discrete

- (i) The number of students in a class.
- (ii) The number of workers in a company
- (iii) The number of parts damaged in during transportation.
- (iv) Shoe sizes.
- (v) Number of languages an individual speaks.
- (vi) The number of test questions you answered correctly.

- The amount of time required to complete a project.
- The height of children.
- The amount of rain, in inches, that falls in a storm.
- The square footage of a two-bedroom house.
- The weight of a truck.
- The speed of light

SET: Group of things that belong together.

⑤ CO's Course Outcomes.

CO1:- Write an argument using logical notation and determine if the argument is or is not valid.

CO2:- Understand the basic principles of sets and operations in sets

CO3:- Demonstrate an understanding of relations and functions and be able to determine their properties.

CO4:- Demonstrate the different traversal methods for trees and graphs.

CO5:- Model problems in Computer Science using graphs and trees.

⑥ Syllabus

Theory (5 units)

Lab. (Programs in C programming and mapple Software)

5 units

UNIT 1:- Set Theory, Functions, Natural Numbers.

UNIT 2:- Algebraic Structures

UNIT 3:- Lattices, Boolean Algebra.

UNIT 4:- Propositional Logic, Predicate Logic.

UNIT 5:- Trees, graphs and Combinatorics.

Planning of Lecture:-

- ⇒ Greetings to Student. (With + Self Name + Subject Name)
- ⇒ Literal meaning of Sets
- ⇒ Student's Introduction.
- ⇒ Motivation to move forward, Now.
- ⇒ Literal meaning of Subject. + why study + COs + syllabus
- ⇒ SET

Common
Question

Student's Introduction.
 Your Name:-
 Meaning of Name:-
 Who are you?
 How are you?
 What is the significance of your name in your life?

UNIT I

Set Theory :- Introduction, Combination of sets, Multisets, ordered pairs. Proofs of some general identities on sets

Relations:- Definition, operations on relations, properties of relations, composite relations, Equality of relations, recursive definition of relation, order of Relations.

Functions:- Definition, classification of functions, operations on functions, recursively defined functions, Growth of functions.

Natural Numbers:- Introduction, Mathematical Induction, Variants of induction, Induction with Nonzero Base Case, proof methods, Proof by Counter-example, Proof by Contradiction.

Lecture 1.IntroductionSETS and Elements

(i) A set is a collection of objects, the elements or members of the set.

Sets denotation:- A, B, X, Y, \dots (capital letters)

elements " :- a, b, x, y, \dots (small letters)
lowercase

~~(ii)~~ $b \in A$:: 'b belongs to A' or 'b is an element of A'
~~(iii)~~ $b \notin A$:: 'b does not belong to A' or 'b is not an element of A'

(iv) Principle of extension:- 'Two sets A and B are equal if they have the same members.'

Ex. (i) $A = \{1, 2, 4, 7\}, B = \{4, 2, 1, 7\}$

$$A = B \checkmark$$

(ii) $A = \{1, 2, 4, 7\}, B = \{1, 2, 3, 7\}$

$$A \neq B$$

(iii)

Specifying Sets :- Two ways.

$$A = \{a, e, i, o, u\}$$

A set containing elements.

$$\neg [b \notin A, e \in A]$$

As a letter in the English alphabet, x is a vowel.

State those properties which characterized the elements in the set.

Ex: $B = \{x : x \text{ is an even integer}, x > 0\}$

\Rightarrow "B is the set of x such that x is an even integer and x is greater than 0". \rightarrow denotes the set 'B' whose elements are positive integers.

Ex: $E = \{x : x^2 - 3x + 2 = 0\}$

$\rightarrow E$ is a solution set of given equation.

Ex: $F = \{2, 1\}$

$\therefore E = \{1, 2\}$

$$G = \{1, 2, 2, 1, \frac{6}{3}\} \text{, then } E = F = G.$$

[A set remains the same if its elements are repeated or rearranged] property

Standard Notations.

N = the set of positive integers

Z = the set of integers

Q = the set of rational numbers $\left[\frac{a}{b} \right]$

R = the set of real numbers $\left[a, b \text{ are integers} \right]$

C = the set of complex numbers $\left[2+3i \right]$

(iv) Principle of Abstraction: - "Given any set U and any property P , there is a set A such that the elements of A are exactly those members of U which have the property P ." \rightarrow [describing a set.]

(v) Universal set and Empty set.

→ In any application of the theory of sets, the members of all sets under investigation usually belong to some fixed large set called the universal set:

for ex:- In plane geometry, the universal set consists of all the points in the ~~one~~ plane.

Ex:- In Human population studies the universal set consists of all people in the world.

Symbol: \boxed{U}

Empty set:- For a given set U and a property P , there may not be any elements of U which have property P . For e.g.:

Ex:- $S = \{x : x \text{ is a positive integer, } x^2 = 3\}$
has no elements since no positive integer has the required property.

⇒ The set with no elements is called the empty set, or null set and is denoted by \emptyset .

$S = \{\}$ → Set representation.
 $S = \emptyset$

$T = \{\emptyset\}$ → A set is having an empty set.

$S \neq T$

(vi) SUBSETS:- If every element of a set A is also an element of a set B , then A is called a subset of B .

$A \subseteq B$ or $B \supseteq A$

" A contained in B " or " B contain A ".

$A \not\subseteq B$ or $B \not\supseteq A \Rightarrow A$ is not a subset of B .

\Rightarrow at least one element of A does not belong to B .

Ex:- $A = \{1, 3, 4, 5, 6, 9\}$, $B = \{1, 2, 3, 5, 7\}$, $C = \{1, 5\}$

Relation

$\boxed{N \subseteq Z \subseteq Q \subseteq R}$

$$\rightarrow E = \{2, 4, 6\}, \quad F = \{6, 2, 4\}$$

(4) 6

$$E \subseteq F \text{ and } F \subseteq E$$

$$\Rightarrow E = F \quad \checkmark$$

* \Rightarrow Every set is a subset of itself. ✓ (since, trivially, elements of A belong to every element of A).

* \Rightarrow Every set A is a subset of the universal set U since since all the elements of A belong to U.

\Rightarrow The empty set \emptyset is a subset of A.

$$\{\}$$

$$\{__, \}$$

$A \subseteq B$ { every element of A belongs to B}
 $B \subseteq C$ { every element of B belongs to C}

$\Rightarrow A \subseteq C$, {since every element of A belongs to C}

Theorem (i) for any set A, we have $\emptyset \subseteq A \subseteq U$

(ii) For any set A, we have $A \subseteq A$

(iii) If $A \subseteq B$ and $B \subseteq C$, then $A \subseteq C$

(iv) $A = B$ if and only if $A \subseteq B$ and $B \subseteq A$

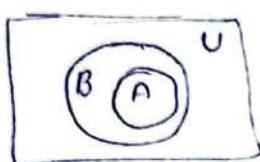
* \Rightarrow If $A \subseteq B$, then it is still possible that $A = B$. When $A \subseteq B$ then but $A \neq B$, we say A is a proper subset of B.

$\Rightarrow A \subset B$

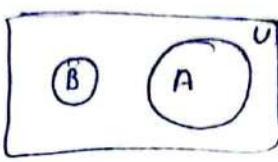
$$\text{Ex: } A = \{1, 3\} \quad B = \{1, 2, 3\}, \quad C = \{1, 2, 3\}$$

\Rightarrow A and B are both subsets of C; but A is a proper subset of C, whereas B is not a proper subset of C since $B = C$.

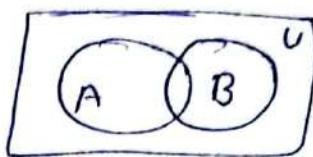
VENN DIAGRAM :- Pictorial representation of sets in which sets represented by enclosed area in the plane. (5)



(a) A ⊂ B



(b) A and B are Disjoint



(c)

→ Some objects are in A, but not in B.

→ Some are in B but not in A.

→ Some are in both

→ Some are in neither A nor B

Arguments and Venn diagrams

Ex:-

S1: My saucepans are the only things I have that are made of tin.

S2: I find all your presents very useful.

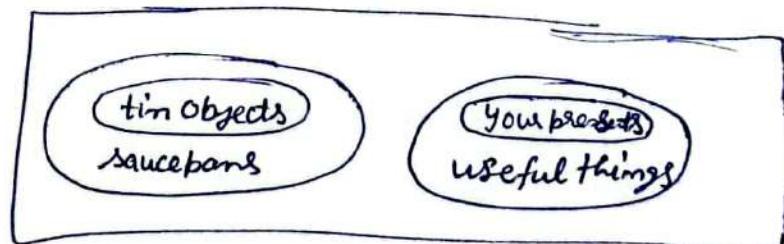
S3: None of my saucepans is of the slightest use.

S : Your presents to me are not made of tin

S₁, S₂, S₃ are assumptions

S \Rightarrow Conclusion

\Rightarrow By S



By S₁; the ~~tin~~ objects are contained in the set of saucepans and by S₃ the set of saucepans and the set of useful things are disjoint;

By S₂ the set of "your presents" is a subset of the set of useful things. hence draw fig.

\Rightarrow The conclusion is clearly valid by the above Venn diagram because the set of "your presents" is disjoint from the set of tin objects.

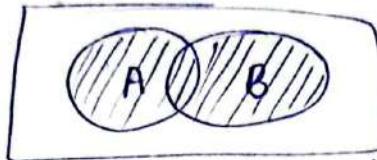
SET OPERATIONS

(1) Union and Intersection

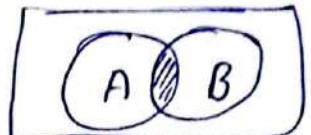
⇒ The union of two sets A and B, denoted by $A \cup B$, is the set of all elements which belong to A or to B; that is

$$A \cup B = \{x : x \in A \text{ or } x \in B\}$$

Here "or" is used in the sense and/or,



(a) $A \cup B$ is shaded



(b) $A \cap B$ is shaded

The intersection of two sets A and B, denoted by $A \cap B$, is the set of elements which belong to both A and B; that is

$$A \cap B = \{x : x \in A \text{ and } x \in B\}$$

Ex:- Let $A = \{1, 2, 3, 4\}$, $B = \{3, 4, 5, 6, 7\}$, $C = \{2, 3, 5, 7\}$

$$A \cup B = \{ \quad \quad \quad \}$$

$$A \cap B = \{ \quad \quad \quad \}$$

$$A \cup C = \{ \quad \quad \quad \}$$

$$A \cap C = \{ \quad \quad \quad \}$$

Ex:- Let J contains the set of AIML-I students.

Let M denotes the set of male students in GL Bajaj, and let F denotes.

$$M \cup F = C$$

$$M \cap F = \emptyset.$$

Theorem: The following are equivalent: $A \subseteq B$, $A \cap B = A$, and $A \cup B = B$

(2) Compliments:

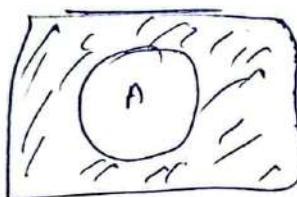
A^c , is the set of elements which belong to U but do not belong to A ; that is

$$A^c = \{x : x \in U, x \notin A\}$$

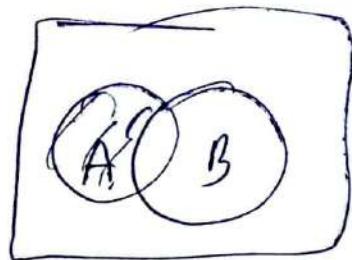
$$\left\{ \begin{array}{l} \text{Other Denotations} \\ A^c = A' = \bar{A} \end{array} \right\}$$

→ The relative complement of a set B with respect to a set A or, simply, the difference of A and B , denoted by $A \setminus B$, is the set of elements which belong to A but which do not belong to B ; that is

$$A \setminus B = \{x : x \in A, x \notin B\}$$



A^c is shaded.



$A \setminus B$ is shaded

Ex:- $U = N = \{1, 2, 3, \dots\}$, the positive integers, is the universal set. $E = \{2, 4, 6, 8, \dots\}$ every even integer.
 $A = \{1, 2, 3, 4\}$, $B = \{3, 4, 5, 6, 7\}$, $C = \{6, 7, 8, 9\}$ $\rightarrow A^c, B^c, C^c$
 $A \setminus B, B \setminus C, B \setminus A, C \setminus E$

(4) Symmetric Difference.

$$A \oplus B = (A \cup B) \setminus (A \cap B)$$

$$A \oplus B = (A \setminus B) \cup (B \setminus A)$$

$$A = \{1, 2, 3, 4, 5, 6\} \text{ and } B = \{4, 5, 6, 7, 8, 9\}$$

(5) Fundamental Products:- Consider n distinct sets A_1, A_2, \dots, A_n . A fundamental product of the sets is a set of the form

$$A_1^* \cap A_2^* \cap \dots \cap A_n^*$$

where A_i^* is either A_i or A_i^c .

(1) There are 2^n such fundamental products.

(2) Any two such fundamental products are disjoint, and

(3) the universal set U is the union of all the fundamental products.

Ex: $U = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$ (8)

$$A = \{1, 2, 3\}$$

$$B = \{6, 7, 8\}$$

$$C = \{4, 5, 9\}$$

Algebra of Sets and DUALITY

sets under the operations of union, intersection, and complement.
satisfy various laws or identities which are listed below.

$$\boxed{(A \cup B)^c = A^c \cap B^c} \rightarrow \text{DeMorgan's laws}$$

(1) Idempotent laws

$$(1a) A \cup A = A$$

~~(2a) Associative law~~
$$(2a) (A \cup B) \cup C = A \cup (B \cup C)$$

~~(3a) Commutative law~~
$$(3a) A \cup B = B \cup A$$

$$(1b) A \cap A = A$$

$$(2b) (A \cap B) \cap C = A \cap (B \cap C)$$

$$(3b) A \cap B = B \cap A$$

(2) Distributive laws.

$$(4a) A \cup (B \cap C) = (A \cup B) \cap (A \cup C) \quad (4b) A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$$

(3) Identity laws

$$(5a) A \cup \emptyset = A$$

$$(6a) A \cup U = U$$

$$(5b) A \cap U = A$$

$$(6b) A \cap \emptyset = \emptyset$$

~~(4) Involution laws~~

$$(7) A^c = (A^c)^c = A$$

Complement Laws.

$$8(a) A \cup \bar{A} = U$$

$$8(b) A \cap \bar{A} = \emptyset$$

$$9(a) U^c = \emptyset$$

$$9(b) \emptyset^c = U$$

DeMorgan's Laws

$$(10a). (A \cup B)^c = A^c \cap B^c$$

$$(10b) (A \cap B)^c = A^c \cup B^c$$

* Principle of Duality.

Suppose E is an equation of set algebra. The dual E^* of E is the equation obtained by replacing each occurrence of \cup , \cap , U , and ϕ in E by \cap , \cup , ϕ , and U respectively.

union universal set \downarrow
 \cap universal set union \downarrow
 \cup universal set \cap universal set

E

for ex - the dual of.

$$(U \cap A) \cup (B \cap A) = A \text{ is } (\phi \cup A) \cap (B \cup A) = A$$

→ This fact is known as principle of duality.
 Concept

If any equation E is an identity, then its dual E^* is also an identity.

key points (i) If A and B are disjoint finite sets, then $A \cup B$ is finite and

$$n(A \cup B) = n(A) + n(B)$$

(ii) If A and B are finite sets, then $A \cup B$ and $A \cap B$ are finite and.

$$n(A \cup B) = n(A) + n(B) - n(A \cap B)$$

(iii) If A , B , and C are finite sets, then so is $A \cup B \cup C$, and

$$n(A \cup B \cup C) = n(A) + n(B) + n(C) - n(A \cap B) - n(A \cap C) - n(B \cap C) + n(A \cap B \cap C).$$

Ex: Consider the following data for 120 Mathematical students at a college concerning the languages French, German, and Russian:

65 study French

45 study German, 42 study Russian.

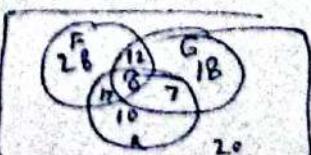
20 study French and German

25 study French and Russian

15 study German and Russian.

8 study all three languages.

How many students study only one of the languages?



Classes of sets, Power sets & Partitions

Multisets

A multiset is an ordered collection of elements, in which the multiplicity of an element may be one or more than one or zero. The multiplicity of an element is the number of times the element repeated in the multiset.

$$A = \{l, l, m, m, n, n, n, n\}$$

$$B = \{a, a, a, a, a, c\}$$

Operations

(i) Union of multisets

$$A = \{l, l, m, m, n, n, n\}$$

$$B = \{l, m, m, m, n\}$$

$$A \cup B = \{l, l, m, m, m, n, n, n, n\}$$

} \Rightarrow the multiplicity of an element is equal to the maximum of the multiplicity of an element in A and B.

(ii) Intersection of multisets

$$A = \{l, l, m, n, p, q, q, r\}$$

$$B = \{l, m, m, p, q, q, r, r, r, r\}$$

$$A \cap B = \{l, m, p, q, r\}$$

} the multiplicity of an element is equal to the minimum of the multiplicity of an element in A and B.

(iii) Difference of multisets

$$A = \{l, m, m, m, n, n, n, p, p, b\}$$

$$B = \{l, m, m, m, n, n, n, r, r, r\}$$

$$A - B = \{n, n, p, p, b\}$$

} the multiplicity of an element is equal to the multiplicity of the element in A minus the multiplicity of the element in B if the difference is +ve, and is equal to zero if the difference is 0 or negative?

Sum of Multisets: The sum of two multisets A and B, is a multiset such that the multiplicity of an element is equal to the sum of the multiplicity of an element in A and B.

$$A = \{l, m, n, p, q\}$$

$$B = \{l, l, m, n, n, n, p, q, q\}$$

$$A+B = \{l, l, l, m, m, n, n, n, n, p, p, q, q, q\}$$

Cardinality of sets - \Rightarrow : the no. of distinct elements in a multiset without considering the multiplicity of an element.

$$A = \{l, l, m, m, n, n, n, p, p, p, p, q, q, q\}$$

Cardinality of set A is 5.

Ordered Set - \Rightarrow ordered collection of distinct objects,

~~It is~~

$$\text{Hall no. } \{3, 6, 7, 8, 9\}$$

$$\text{Week days} = \{S, M, T, W, W, TH, F, S, S\}$$

(a, b) and (b, a) are two different ordered pairs.

$\underbrace{\{(a, b), c\}}_{\text{Int. element or ordered pair}} \rightarrow \text{ordered triple}$ $\underbrace{\{((a, b), c), d\}}_{\text{Int. element or ordered triple}}$

Ordered Quadable.

Ex: ordered set of 5 elements

$$\underbrace{\{(((a, b), c), d)\}}_{(n-1)} e$$

\downarrow
etc

Relations

(12)

Relations will be defined as in terms of ordered pairs (a, b) of elements, where a is designated as first element and b as second element.

$\rightarrow (a, b) = (c, d)$ if and only if $a = c$ and $b = d$.

Then $(a, b) \neq (b, a)$ unless $a = b$.

Product Sets

Two sets A, B . The set of all ordered pairs (a, b)

where $a \in A$ and $b \in B$ is called the product, or Cartesian product, of A and B .

$$A \times B = \{(a, b) : a \in A \text{ and } b \in B\}$$

for $A \times A$, can also be written as A^2 .

Ex: R denotes the real numbers. So $R^2 = R \times R$, is the set of ordered pairs of real numbers.

Ex: $A = \{1, 2\}, B = \{a, b, c\}$ then

$$A \times B = \{(1, a), (1, b), (1, c), (2, a), (2, b), (2, c)\}$$

$$B \times A = \{(a, 1), (a, 2), (b, 1), (b, 2), (c, 1), (c, 2)\}$$

Also $A \times A = \{(1, 1), (1, 2), (2, 1), (2, 2)\}$

$$(A \times B \neq B \times A)$$

$$n(A \times B) = 6 = 2 \cdot 3 = n(A) \cdot n(B)$$

Relations

Def: Let A and B be sets. A binary relation, or, simply, relation from A to B is a subset of $A \times B$.

- (i) $(a, b) \in R$; we then say "a is R -related to b ", written aRb
- (ii) $(a, b) \notin R$; we say "a is not R -related to b ", written $a \not R b$

$$A^2 = A \times A \Rightarrow R \text{ is a relation on } A.$$

* The domain of a relation R is the set of all first elements of the ordered pairs which belong to R , and the range of R is the set of 2nd elements.

Ex:-

$$(a) A = \{1, 2, 3\} \text{ and } B = \{x, y, z\}$$

$$\text{and let } R = \{(1, y), (1, z), (3, y)\}$$

then R is a relation from A to B . Since R is a subset of $A \times B$.

$$1Ry, 1Rz, 3Ry, \text{ but } 1Rx, 2Rx, 3Rz, 3Rx, 2Rz, \\ \cancel{2Ry}, 2Ry.$$

(b) A familiar relation on the set \mathbb{Z} of integers is "division". A common notation for this relation is to write m/n when m divides n . Thus $6|30$ but $7 \nmid 25$.

(c). Consider the set L of lines in the plane. Perpendicularity, written \perp , is a relation on L . That is given any pair of lines a and b , either $a \perp b$, or $a \not R b$.

- (A). Let A be any set. Then $A \times A$ and \emptyset are subsets of $A \times A$ (14)
 hence are relations on A called the universal relation or empty relation.

Inverse Relation:-

- Let R be any relation from a set A to a set B. The inverse of R, denoted by R^{-1} , is the relation from B to A which consists of those ordered pairs, which when reversed, belong to R; that is,

$$R^{-1} = \{ (b, a) : (a, b) \in R \}$$

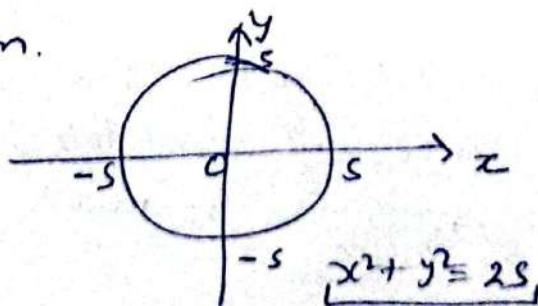
Ex let $R = \{(1, y), (1, z), (3, y)\}$ from $A = \{1, 2, 3\}$
 $B = \{x, y, z\}$

$$R^{-1} = \{(y, 1), (z, 1), (y, 3)\}$$

$$(R^{-1})^{-1} = R$$

Pictorial representation:- $x^2 + y^2 = 2s$,
 relation on S.

- S consists of all ordered pairs (x, y) which satisfy the given equation.

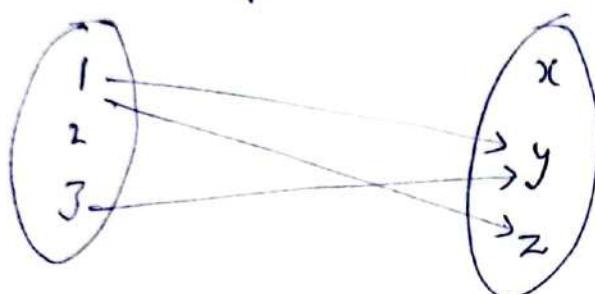


Representation of Relations on finite sets. (Two ways)

(1)

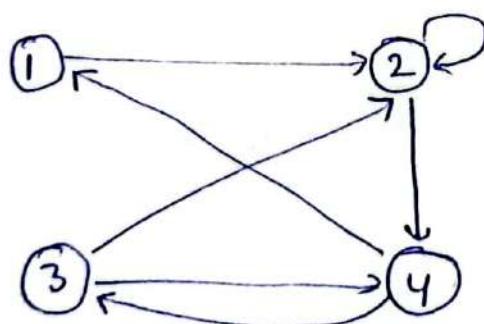
	x	y	z
1	0	1	1
2	0	0	0
3	0	1	0

(ii)



Directed graph of Relations on sets.:

$$R = \{(1,2), (2,2), (2,4), (3,2), (3,4), (4,1), (4,3)\}$$



Composition of relations.

A, B, C are sets, R & S are relations.

$A^R B$, $B^S C$ \Rightarrow R is a subset of $A \times B$.

\Rightarrow S is a subset of $B \times C$

then R and S give rise to a relation from A to C denoted by $R \circ S$ and defined by

$a(R \circ S)_C$ if for some $b \in B$ we have aRb and bSc .

That is.

$R \circ S = \{(a,c) : \text{there exist } b \in B \text{ for which } (a,b) \in R \text{ and } (b,c) \in S\}$

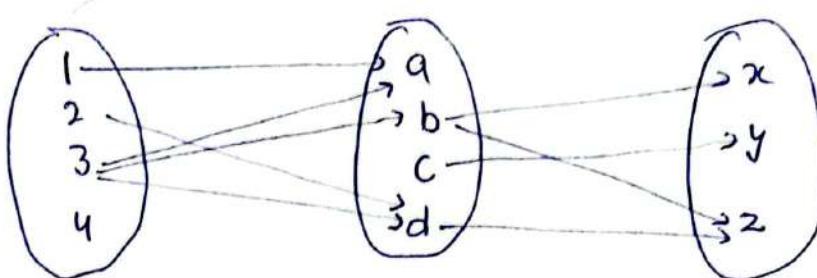
→ Composition of R with itself is always defined. (10)

$$R \circ R \in R^2$$

→ $R^3 = R \circ R \circ R = R^2 \circ R$ and so on. $\underline{R^n}$ is defined for all
+ve n.

Ex: Let $A = \{1, 2, 3, 4\}$, $B = \{a, b, c, d\}$, $C = \{x, y, z\}$ and
let $R = \{(1, a), (2, d), (3, a), (3, b), (3, d)\}$
and $S = \{(b, x), (b, z), (c, y), (d, z)\}$.

↓
Diagram



✓ ~~$(R \circ S)_z$~~ since $2 R_d$ and $d S_z$

✓ ~~$(R \circ S)_x$~~ and ~~3~~

$$R \circ S = \{(2, z), (3, x), (3, z)\} \quad \checkmark$$

Types of Relations.

(i) Reflexive Relations: A relation R on a set A is reflexive if aRa for every $a \in A$, that is, if $(a, a) \in R$ for every $a \in A$. Thus R is not reflexive if there exist an $a \in A$ such that $(a, a) \notin R$.

Ex: ~~Ques.~~ $A = \{1, 2, 3, 4\}$

$$R_1 = \{(1, 1), (1, 2), (2, 3), (1, 3), (4, 4)\}$$

$$R_2 = \{(1, 1), (1, 2), (2, 1), (2, 2), (3, 3), (4, 4)\}$$

$$R_3 = \{(1, 3), (2, 1)\} \quad R_4 = \emptyset, \text{ empty relation}$$

$$R_5 = A \times A, \text{ universal relation}$$

which are
its
reflexive

Lecture

(17)

Example 1

Consider the following relations -

on set $A = \{1, 2, 3, 4\}$

$R_1 = \{(1,1), (1,2), (2,3), (1,3), (4,4)\} \rightarrow \text{X}$

$R_2 = \{(1,1), (1,2), (2,1), (2,2), (3,3), (4,4)\} \rightarrow \checkmark$

$R_3 = \{(1,3), (2,1)\} \rightarrow \text{X}$

$R_4 = \text{empty relation} \rightarrow \text{X}$

$R_5 = A \times A \text{ the universal relation} \rightarrow \checkmark$

Reflexive	Symmetric	Antisym.	Transitive	Equivalence
→ X	→ ✓	→ ✓	→ ✓	→
→ ✓	→ ✗	→ ✗	→ ✓	→
→ X	→ ✓	→ ✓	→ X	→
→ ✓	→ ✓	→ ✓	→ ✓	→
→ ✓	→ ✗	→ ✓	→ ✓	→

Ex 2:- Consider the following five relations.

R

(i) Relation \leq (less than or equal) on the set Z of integers. →

(ii) Set inclusion \subseteq on a collection C of sets →

(iii) Relation \perp (perpendicular) on the set L of lines in the plane. → X

(iv) Relation \parallel (parallel) on the set L of lines in the plane → X

(v) Relation $|$ of divisibility on the N of positive integers : →

Ref.	Sym.	Antisym.	Transitivity	Equivalence
✓	→ X	→ ✓	→ ✓	→
✓	→ X	→ ✓	→ ✓	→
X	→ ✓	→ X	→ X	→ X
X	→ ✓	→ X	→ X	→ X
✓	→ X	→ ✓	→ ✓	→ ✓

Types of Relations.

(A) Reflexive:- A relation R on a set A is reflexive if aRa for every $a \in A$, that is, if $(a,a) \in R$ for every $a \in A$. Thus R is not reflexive if there exists an $a \in A$ such that $(a,a) \notin R$.

(B) Symmetric:- A relation R on a set A, is symmetric if whenever aRb then bRa , that is, if whenever $(a,b) \in R$ then $(b,a) \in R$. Thus R is not symmetric if there exists ~~(a,b)~~ $a, b \in A$ such that $(a,b) \in R$ but $(b,a) \notin R$.

≡

(C) Antisymmetric:- A relation R on a set A is antisymmetric if + whenever aR_b and bR_a then $a=b$, that is, if whenever $(a,b), (b,a) \in R$ then $a=b$. Thus R is not antisymmetric if there exists $(a,b) \in R$ such that (a,b) and (b,a) belong to R , but $a \neq b$.

Remark:- The properties of being symmetric and being antisymmetric are not negatives of each other. For example, the relation R $R = \{(1,3), (3,1), (2,3)\}$ is neither symmetric nor antisymmetric. On the other hand, the relation $R' = \{(1,1), (2,2)\}$ is both symmetric and antisymmetric.

(D) Transitive:- A relation R on a set A is transitive if whenever aR_b and bR_c then aR_c , that is, If whenever $(a,b), (b,c) \in R$ then $(a,c) \in R$. Thus R is not transitive if there exists $a, b, c \in A$ such that $(a,b), (b,c) \in R$ but $(a,c) \notin R$.

$\Rightarrow \perp$ is not transitive. If $a \perp b$ and $b \perp c$, then it is not true that $a \perp c$. Since no line is parallel to itself, we can have $a \parallel b$ and $b \parallel c$, but $a \not\parallel c$. Thus \parallel is not transitive.

(E) Equivalence Relations:- Consider a nonempty set A . A relation R on A is an equivalence relation if R is reflexive, symmetric and transitive. That is, R is an equivalence relation on A if it has the following three properties.

- (i) for every $a \in A$
- (ii) if $a R b$, then $b R a$
- (iii) if $a R b$ and $b R c$, then $a R c$

$| - |$

$(a,b), (b,a)$
I. I. I

for exple. alike
objects.

- (i) If $a = a$ for all $a \in A$
- (ii) If $a = b$, then $b = a$
- (iii) If $a = b$, and $b = c$,
then $a = c$.

closure Properties.

(17)

Reflexive and Symmetric closure

→ here $\Delta_A = \{(a,a), a \in A\}$ is the diagonal or equality relation on A.

Th: Let R be a relation on set A. Then:

(i) $R \cup \Delta_A$ is the reflexive closure of R.

(ii) $R \cup R^{-1}$ is the symmetric closure of R.

In other words, reflexive(R) is obtained by simply adding to R those elements (a,a) in the diagonal which do not already belong to R, and symmetric(R) is obtained by adding to R all pairs (b,a) whenever belong to R.

Ex: $A = \{1, 2, 3, 4\}$

$$R = \{(1,1), (1,3), (2,4), (3,1), (3,3), (4,3)\}$$

Then, $\text{reflexive}(R) = R \cup \{(2,2), (4,4)\}$ and

$$\text{symmetric}(R) = R \cup \{(4,2), (3,4)\}$$

Ex: Consider the relation $<$ (less than) on the set N of positive integers. Then

$$\text{reflexive}(<) = < \cup \Delta = \leq = \{(a,b); a \leq b\}$$

$$\text{symmetric}(<) = < \cup > = \{(a,b); a \neq b\}$$

Transitive closure:-

Let R be a relation on a set A. Recall that $R^2 = R \circ R$ and

$R^n = R^{n-1} \circ R$, we define

$$R^\ast = \bigcup_{i=1}^{\infty} R^i$$

The following thm applies,

Th 2.4: R^\ast is the transitive closure of a relation R.

Suppose A is a finite set with n elements. Then we show that

$$R^\ast = R \cup R^2 \cup R^3 \cup \dots \cup R^n$$

So, $\text{transitive}(R) = R \cup R^2 \cup \dots \cup R^n$

Ex: R on $A = \{1, 2, 3\}$

$$R = \{(1,2), (2,3), (3,1)\}$$

Then $R^2 = R \circ R = \{(1,3), (2,3), (3,1)\}$ and

$$R^3 = R^2 \circ R = \{(1,3), (2,3), (3,1)\}$$

accordly

$$\text{transitive}(R) = R \cup R^2 \cup R^3 = \{(1,2), (2,3), (3,1)\}$$

Equivalence Relations and Partitions: :-

\Rightarrow A Partition P of S is collection $\{A_i\}$ of nonempty ~~sets~~ subsets of S with the following two properties:

(1) Each $a \in S$ belongs to some A_i

(2) If $A_i \neq A_j$, then $A_i \cap A_j = \emptyset$

In other words, a partition P of S is a subdivision of S into disjoint nonempty sets.

\Rightarrow Suppose R is an equivalence relation on a set S . ~~to different nonempty sets~~. For each a in S , let $[a]$ denote the set of elements of S ~~to~~ to which a is related under R ; that is

$$[a] = \{x : (a, x) \in R\}$$

Partial ordering Relation: - A relation R on a set S is called a partial ordering or a partial order if R is reflexive, antisymmetric, and transitive. A set S with a partial ordering R is called a partially ordered set or poset.

Ex(1) \leq The relation \subseteq of set inclusion is a partial ordering on the set ~~N~~ of positive integers on any collection of sets since set inclusion has the three desired properties that is,

- (i) $A \subseteq A$ for any set A (ii) If $A \subseteq B$ and $B \subseteq A$ then $A = B$
Ex(2) " divides" (iii) If $A \subseteq B$ and $B \subseteq C$, then $A \subseteq C$

Equivalence Relation Example

(21)

(a) Consider the set L of lines and the set T of triangles in the Euclidean plane. The relation "is parallel to or identical" is an equivalence relation on L , and congruence and similarity are ~~equivalent~~ equivalence relations on T .

Congruent meaning in Maths :- Two shapes are congruent if they have the same shape and size. We can also say if two shapes are congruent, then the mirror image of one shape is the same as the other.

(b). The classification of animals by species, that is, the relation "is of the same species as", is ~~not~~ an equivalence relation on the set of animals.

✓ (c). The relation \subseteq of set of inclusion is not an equivalence relation. It is reflexive and transitive, but it is not symmetric since $A \subseteq B$ does not imply $B \subseteq A$.

✓ (d) Let m be a fixed positive integer. Two integers a and b are said to be congruent modulo m , written,

$$a \equiv b \pmod{m}$$

if m divides $a-b$.

For example, $m=4$ we have $11 \equiv 3 \pmod{4}$ since 4 divides $11-3$, and $22 \equiv 6 \pmod{4}$ since 4 divides $22-6$.

⇒ The relation of Congruence of modulo m is an equivalence relation.

Note:-

⇒ The property of transitivity can also be expressed in terms of the composition of relations. For a relation R on A we define

$$R^2 = R \circ R \text{ and, more generally, } R^n = R^{n-1} \circ R$$

Then we have the following result.

Th:- A relation R is transitive if and only if $R^n \subseteq R$ for $n \geq 1$

(A) Partial Order:- A binary relation R on a set A is a partial (22)
order if and only if it is
(1) reflexive
(2) antisymmetric, and
(3) transitive.

→ The ordered pair $\langle A, R \rangle$ is called a poset (partially ordered set) when R is a partial order.

Ex:- The less-than-or-equal-to relation on the set of integers I is a partial order, and the set I with this relation is a poset.

(B) Total order:- A binary relation R on a set A is a total order if and only if it is
(1) a partial order, and
(2) for any pair of elements a and b of A , $\langle a, b \rangle \in R$ or $\langle b, a \rangle \in R$.

⇒ That is, every element is related with every element one way or the other.

A total order is also called a linear order.

Ex:- The less-than-or-equal-to relation on the set of integers I is a total order.

→ The strictly-less-than and proper-subset relations are not partial order because they are not reflexive. They are examples of some relation called quasi order.

(C) Quasi order:- A binary relation R on a set A is a quasi order if and only if it is (No element is related to itself)
(1) irreflexive, and
(2) transitive.

⇒ A Quasi order is necessarily antisymmetric as one can easily verify.

Ex:- (i) less than
(ii) proper subset -

CHAPTER 3

(23)

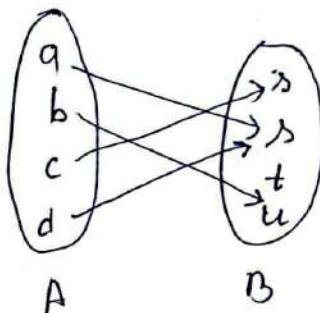
Functions

Defⁿ: Suppose that to each element of a set A we assign a unique element of a set B; the collection of such assignment is known called a function from A into B. The set A is called the domain of the function, and the set B is called the codomain.

$$\underbrace{f: A \rightarrow B}_{\downarrow}$$

f is a function from A into B.

Ex: $f(x) = x^2$ or $x \rightarrow x^2$ or $y = x^2$
 ↓ ↓ ↓
 function variable "goes into" y is dependent variable
 x is independent variable.



Identity function: Let A be any set. The function from A into A which assigns to each element that element itself is called the identity function on A and is usually denoted by I_A , or simply I. In other words.

$$I_A(a) = a \text{ for every element } a \text{ in } A.$$

Defⁿ: A function $f: A \rightarrow B$ is a relation from A to B (i.e., a subset of $A \times B$) such that each $a \in A$ belongs to a unique ordered pair (a, b) in f.

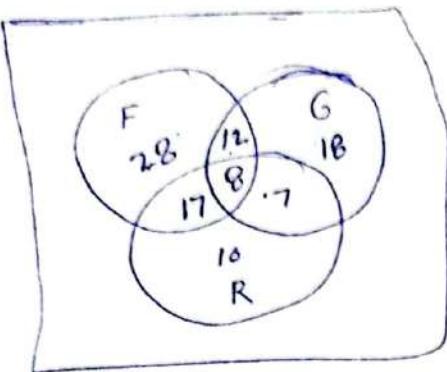
Ex: Consider the following relations on the set $A = \{1, 2, 3\}$:

$$f = \{(1, 3), (2, 3), (3, 1)\}$$

$$g = \{(1, 2), (3, 1)\}$$

$$h = \{(1, 3), (2, 1), (1, 2), (3, 1)\}$$

} which are its not function



Total 120 Students

CS, AI, ML

(Hindi, English, Urdu)

$$\begin{aligned}
 n(A \cup B \cup C) &= n(A) + n(B) + n(C) - n(A \cap B) - n(A \cap C) - n(B \cap C) + \\
 &\quad n(A \cap B \cap C) \\
 &= 65 + 45 + 42 - 20 - 25 - 15 + 8 \\
 &= 100
 \end{aligned}$$

120 - 100 = 20 Students ~~do not study any of the language~~

$20 - 8 = 12$ study F.

$$\underline{\underline{28 + 18 + 10 = 56}}$$

Let there are 120 students in Language College and three languages to study. Hindi, Urdu and Sanskrit.

65 study Hindi

45 study Urdu

42 study Sanskrit

20 study Hindi and Urdu

25 study Hindi and Sanskrit.

15 study Urdu and Sanskrit

8 study all three languages.

How many students are there who study only one language?

Composition functions:-

24 (8)

Let $f: A \rightarrow B$ and $g: B \rightarrow C$; that is, where Codomain of f is the domain of g . Then we may define a new function from A to C , called the composition of f and g and written gof .

$$(gof)(a) = g(f(a)),$$

Note:- Consider any function $f: A \rightarrow B$, then

$$f \circ I_A = f \text{ and } I_B \circ f = f$$

where I_A and I_B are the Identity functions on A and B , respectively.

One to one functions:-

A function $f: A \rightarrow A$ is said to be one-to-one if different elements in the domain A have distinct images.

or

f is one to one if $f(a) = f(a')$ implies $a = a'$

Onto functions:-

A function $f: A \rightarrow B$ is called said to be an onto function if each element of B is the image of some element of A .

In other words, $f: A \rightarrow B$ is onto if the image of f is the entire codomain, i.e., if $f(A) = B$.

$\Rightarrow f$ is a function from A onto B .

f maps A onto B

Invertible function:- A function $f: A \rightarrow B$ is invertible if its inverse relation f^{-1} is a function from B to A .

Thm:- A function $f: A \rightarrow B$ is invertible if and only if f is both one-to-one and onto.

Floor and Ceiling functions

(25) (20)

$\lfloor x \rfloor$, called the floor of x , denotes the greater integer that does not exceed x .

$\lceil x \rceil$, called the ceiling of x , denotes the least integer that is not less than x .

If x is itself an integer, $\lfloor x \rfloor = \lceil x \rceil$, otherwise $\lfloor x \rfloor + 1 = \lceil x \rceil$

$$\lfloor 3.14 \rfloor = 3 \quad \lfloor \sqrt{5} \rfloor = 2 \quad \lfloor -8.5 \rfloor = -9 \quad \lceil -8.5 \rceil = -8$$

$$\lfloor 7 \rfloor = 7, \quad \lceil 7 \rceil = 7.$$

Integer and Absolute Value function

→ The integer value of x , written $\text{INT}(x)$, converts x into an integer by deleting (truncating) the fractional part of the number. Thus,

$$\text{INT}(3.14) = 3 \quad \text{INT}(\sqrt{5}) = 2 \quad \text{INT}(-8.5) = -8, \quad \text{INT}(7) = 7$$

→ The absolute value of the real number x , written $\text{ABS}(x)$ or $|x|$, is defined as the greater of x or $-x$.

$$|-15| = 15, \quad |7| = 7, \quad |-3.33| = 3.33, \quad |4.44| = 4.44.$$

⇒ $|x| = |-x|$ and, for $x \neq 0$, $|x|$ is positive.

Recursively defined functions :-

A function is said to be recursively defined ~~function~~ if the function definition refers to itself.

Two essential properties,

- 1). There must be certain arguments, called base values, for which the function does not refer to itself.
- 2). Each time the function does not refer to itself, the argument of the function must be closer to a base value.

Ex: factorial function, fibonacci sequence.

Growth of functionAlgorithms and Functions $f(x)$

→ An algorithm M is a finite step by step list of well-defined instructions for solving a particular problem say, to find the output $f(x)$ for a given function f with input x . There may be more than one way to obtain $f(x)$, The particular choice of the algorithm M to obtain $f(x)$ may depend on the "efficiency" or "complexity" of the algorithm.

Complexity of Algorithm.

- The Complexity of an algorithm M is the function $f(n)$ which gives the running time and/or storage space requirement of the algorithm in terms of the size n of the input data.

Two Cases

(i) Worst case :- The maximum value of $f(n)$ for any possible input.

(ii) Average case :- The expected value of $f(n)$.

Rate of Growth:- Big O Notation

Suppose M is an algorithm, and suppose n is the size of the input data. Clearly the complexity $f(n)$ of M increases as n increase.

It is usually the rate of increase of $f(n)$ that we want to examine.

$$\log_2 n, n, n \log_2 n, n^2, n^3, 2^n.$$

Defn: Let $f(x)$ and $g(x)$ be arbitrary functions defined on R or a subset of R , we say " $f(x)$ is of order $g(x)$ ", written,

$$f(x) = O(g(x))$$

If there exists a real number k and a positive constant C such that, for all $x > k$, we have

$$\bullet |f(x)| \leq C|g(x)|$$

Ex:-

$$\begin{aligned} 7x^2 - 9x + 4 &= O(x^2) \\ 8x^3 + 832x^2 - 248 &= O(x^3) \end{aligned}$$

- ~~Worst case time complexity of well known algorithms:~~
- Linear search: $O(n)$
 - Binary search: $O(\log n)$
 - Bubble Sort: $O(n^2)$
 - merge sort: $O(n \log n)$

Recursive functions:

Let a and b be positive integers, and suppose Q is defined recursively as follows:

$$Q(a, b) = \begin{cases} 0 & \text{if } a < b \\ Q(a-b, b) + 1 & \text{if } b \leq a \end{cases}$$

(a) find (i) $Q(2, 5)$, (ii) $Q(12, 5)$

(b) what does this function Q do? find $Q(5861, 7)$.

Ans (a) (i) $Q(2, 5) = 0$ since $2 < 5$

$$\begin{aligned} \text{(ii)} \quad Q(12, 5) &= Q(7, 5) + 1 \\ &= [Q(2, 5) + 1] + 1 = Q(2, 5) + 2 \\ &= 0 + 2 = 2 \end{aligned}$$

(b) Each time b is subtracted from a , the value of Q is increased by 1. Hence $Q(a, b)$ finds the quotient when a is divided by b . Thus $Q(5861, 7) = 837$.

$$\begin{array}{r} 14,337 \\ \boxed{14,325} \\ \hline \end{array}$$

~~14,337
14,325

12~~

Mathematical Induction

Natural Numbers: 1, 2, 3, 4, ...
 (+ve integers)

Principle of Mathematical Induction

Suppose there is a given statement $P(n)$ involving the natural number n such that

- (i) The statement is true for $n=1$, i.e. $P(1)$ is true, and
- (ii) If the statement is true for $n=k$ (where k is some positive integer), then the statement is also true for $n=k+1$, i.e., truth of $P(k)$ implies the truth of $P(k+1)$.

Then, $P(n)$ is true for all natural numbers n .

Property (i) is simply a statement of fact. There may be situations when a statement is true for all $n \geq 4$. In this case, step I will start from $n=4$ and we shall verify the result for $n=4$. i.e. $P(4)$.

Property (ii) is a conditional property. It does not assert that the given statement is true for $n=k$, but only that if it is true for $n=k$, then it is also true for $n=k+1$. So, to prove that the property holds, only prove the conditional proposition.

→ If the statement is true for $n=k$, then it is also true for $n=k+1$. ↓ Inductive hypothesis.

→ This is routine referred to as Inductive Step.

Ex: $1 = 1^2 = 1$
 $4 = 2^2 = 1 + 3$
 $9 = 3^2 = 1 + 3 + 5$
 $16 = 4^2 = 1 + 3 + 5 + 7$ etc.

It is worth to be noted that the sum of the first two odd natural numbers is the square of second natural number, sum of the first three natural numbers is the square of third natural number and so on. Thus, from this pattern it appear that

$$1+3+5+7+\dots+(2n-1)=n^2 \text{ i.e.}$$

the sum of the first n odd natural numbers is the square of n .

Let us write.

$$\boxed{P(n) : 1+3+5+7+\dots+(2n-1)=n^2}$$

We wish to prove that $P(n)$ is true for all n .

The first step.

$P(1)$ is true:-

$$1=1^2, \text{ i.e. } P(1) \text{ is true.}$$

Inductive step: let $P(k)$ is true for some positive integer k and we need to prove that $P(k+1)$ is true. Since $P(k)$ is true.

$$1+3+5+7+\dots+(2k-1)=k^2$$

Consider

$$\begin{aligned} 1+3+5+7+\dots+(2k-1)+(2(k+1)-1) \\ = k^2 + (2k+1) = (k+1)^2 \end{aligned}$$

Therefore $P(k+1)$ is true and the inductive proof is complete.
Hence, $P(n)$ is true for all natural numbers. Complete

Ex(1) for all $n \geq 1$, prove that

$$1^2+2^2+3^2+4^2+\dots+n^2 = \frac{n(n+1)(2n+1)}{6}$$

Sol:

$$P(n) : 1^2+2^2+3^2+4^2+\dots+n^2 = \frac{n(n+1)(2n+1)}{6}$$

$$P(1) \text{ is true} = \frac{1(1+1)(2 \times 1 + 1)}{6} = \frac{1 \times 3}{6} = 1, \text{ which is true}$$

Let for k is true.

$$1^2 + 2^2 + 3^2 + 4^2 + \dots + k^2 = \frac{k(k+1)(2k+1)}{6}$$

We shall now prove that $P(k+1)$ is also true; Now, we have

$$(1^2 + 2^2 + 3^2 + 4^2 + \dots + k^2) + (k+1)^2$$

$$= \frac{k(k+1)(2k+1)}{6} + (k+1)^2$$

$$= \frac{k(k+1)(2k+1) + 6(k+1)^2}{6}$$

$$= \frac{(k+1)(2k^2 + 7k + 6)}{6}$$

$$= \frac{(k+1)(k+1+1)(2(k+1)+1)}{6}$$

Thus, $P(k+1)$ is true whenever $P(k)$ is true.

Hence, from the principle of mathematical induction, the statement $P(n)$ is true for all natural numbers n .

Ex: (i) Prove that $2^n > n$ for all positive integers n . \rightarrow

(ii) for all $n \geq 1$, prove that

$$\frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \frac{1}{3 \cdot 4} + \dots + \frac{1}{n(n+1)} = \frac{n}{n+1}$$

$$\left\{ \begin{array}{l} 2^k > k \\ \text{metibly 2 b/w} \\ 2 \cdot 2^k > 2k \\ 2^{k+1} > 2k = k+k > k+1 \\ k+1 > k+1 \end{array} \right\}$$

(iii) for every positive integer n , prove that $7^n - 3^n$ is divisible by 4.

(iv) Prove that $(1+x)^n \geq (1+nx)$, for all natural number n , when $x > -1$.

(v) Prove that:

$$1^2 + 2^2 + \dots + n^2 \geq \frac{n^3}{3}, n \in \mathbb{N}$$

(v) Prove the rule of exponent $(ab)^n = a^n b^n$

by using principle of mathematical induction for every natural number.

Sdn (i)

$$\frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \frac{1}{3 \cdot 4} + \dots + \frac{1}{k(k+1)} = \frac{k}{k+1}$$

$$\frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \frac{1}{3 \cdot 4} + \dots + \frac{1}{k(k+1)} + \frac{1}{(k+1)(k+2)}$$

$$= \frac{k}{(k+1)} + \frac{1}{(k+1)(k+2)} = \frac{k(k+2)+1}{(k+1)(k+2)} = \frac{k^2+2k+1}{(k+1)(k+2)}$$

$$\frac{(k+1)^2}{(k+1)(k+2)} = \frac{k+1}{k+2} = \frac{k+1}{(k+1)+1} \quad \underline{\text{Proved}}$$

Ans :-

$7^k - 3^k$ is divisible by 4

We can write $7^k - 3^k = 4d$, where $d \in N$.

$$\begin{aligned} 7^{(k+1)} - 3^{(k+1)} &= 7^{(k+1)} - 7 \cdot 3^k + 7 \cdot 3^k - 3^{(k+1)} \\ &= 7(7^k - 3^k) + (7-3) \cdot 3^k = 7(4d) + (7-3)3^k \\ &= 7(4d) + 4 \cdot 3^k = 4(\underline{7d + 3^k}) \\ &\quad \underline{\text{divisible by 4.}} \end{aligned}$$

Proof Techniques

- * The following three definitions are central to the execution of our proofs.

Definition 1: An integer number n is even if and only if there exists a number k such that $n = 2k$.

Definition 2: An integer number n is odd if and only if there exists a number k such that $n = 2k+1$.

Definition 3: Two integers a and b are consecutive if and only if $b = a+1$.

(A) Direct Proof [Proof by Construction]

→ In a constructive proof one attempts to demonstrate $P \Rightarrow Q$ directly.

→ Simplest and easiest method of proof.

1. Assume that P is true
2. Use P to show that Q must be true.

Theorem 1:- If a and b are consecutive integers, then the sum $a+b$ is odd.

Proof:- Assume that a and b are consecutive integers.

Because a and b are consecutive we know that $b=a+1$. Thus, the sum $a+b$ may be re-written as $2a+1$. Thus, there exists a number k such that $a+b=2k+1$ so the sum $a+b$ is odd.

(B) Proof by Contradiction

→ We can use this to demonstrate $P \Rightarrow Q$ by assuming both P and $\neg Q$ are simultaneous true and deriving a contradiction. When we derive this contradiction it means that one of our

(2)

assumptions was ~~undesirable~~ untenable. Presumably we have either assumed or already proved P to be true so that finding a contradiction implies that $\neg Q$ must be false.

The method of proof by contradiction

1. Assume that P is true
2. Assume that $\neg Q$ is true
3. Use P and $\neg Q$ to demonstrate a contradiction

Theorem 2. If a and b are consecutive integers, then the sum a+b is odd.

Proof: Assume that a and b are consecutive integers. Assume also that the sum a+b is not odd. Because the sum a+b is not odd, there exists no number k such that $a+b=2k+1$. However, the integers a and b are consecutive integers, so we may write the sum a+b as $2a+1$. Thus, we have derived that $a+b \neq 2k+1$, for any integer k and also that $a+b=2a+1$. This is a contradiction. If we hold a and b are consecutive then we know that the sum a+b must be odd.

(C) Proof by Induction

1. Show that a propositional form $P(x)$ is true for some basic case.
2. Assume that $P(n)$ is true for some n, and show that then implies that $P(n+1)$ is true.
3. Then, by the principle of induction, the propositional form $P(x)$ is true for all n greater or equal to the basic case.

Theorem 3. If a and b are consecutive integers, then the sum $\underline{a+b}$ is odd. (34)

Proof: Define the propositional form $F(x)$ to be true when the sum of x and its successor is odd.

(Step 1) Consider the proposition $F(1)$, The sum $1+2=3$ is odd because we can demonstrate there exists an integer k such that $2k+1=3$. Thus; $F(x)$ is true when $x=1$.

(Step 2): Assume that $F(x)$ is true for some x . Then for some x we have that $x+(x+1)$ is odd.

→ we add 1 to both x and $x+1$ which gives the sum $(x+1)+(x+2)$.

→ we claim two things: first, the sum $(x+1)+(x+2)=F(x+1)$. Second, we claim that adding two to any integers does not change that integer's evenness or oddness, with these two observations we claim that $F(x)$ is odd implies $F(x+1)$ is odd.

(Step 3) By the principle of Mathematical induction we thus claim $F(x)$ is odd for all integers x . Thus, the sum of any two consecutive integers is odd.

D Proof by Contrapositive.

Two proposition are equivalent mean that if one can prove $P \Rightarrow Q$ then they have also proved $\neg Q \Rightarrow \neg P$. and vice versa.

Theorem 4. If the sum $a+b$ is not odd, then a and b are not consecutive integers.

Proof:- Assume that the sum of the integers ~~a~~ 35 ~~b~~ 27 on ~~a~~ and ~~b~~ is odd. Then, there exists no integer k such that $a+b = 2k+1$. Thus, $a+b \neq k+(k+1)$ for all integers k . Because $k+1$ is the successor of k , this implies that a and b cannot be consecutive integers. ✓